

# Introduction to Computer Networks



## Foundation

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# Outline

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- **Applications**
- **Network Connectivity**
- **Network Architecture**
- **Network Performance**

# Applications

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- **Most people know about the Internet (a computer network) through applications**
  - **World Wide Web**
  - **On line games**
  - **Email (Gmail, hotmail,...)**
  - **Online Social Network (Facebook, twitter,...)**
  - **Streaming Audio Video (Youtube, ppstream, kkbox, ...)**
  - **File Sharing (dropbox, ...)**
  - **Instant Messaging (Skype, IM+, MSN, Line, WeChat,...)**
  - **...**

# Example of an application

**SHARE ROOM** ShareRoom展示會議

狀態 會議資訊 設置版面 會議管理 請求發言 錄影 主版面 影音 聊天

主版面 桌面分享 白板 文件

- ~~It is not practical to compute  $T(x)$  if  $m$  exceeds 4 or 5.~~
- Based on the state diagram of a code, there is an efficient way [6] to find the free distance for large  $m$  ( up to about 20).
- No general solution to the problem of finding  $d_{free}$  for large values of  $m$  has yet been discovered.
- Some of good codes are referred in Table 1.1 .
- For Example,  $g^{(1)} = 54 = [1\ 0\ 1\ 1\ 0\ 0]$ ,  $g^{(2)} = 64 = [1\ 1\ 0\ 1\ 0\ 0]$

This is my meeting!

影音 2 x 2(default)

ChiaAn 會員1

黃小月 ATO

聊天

ChiaAn, 17:47  
Hello world

ChiaAn, 17:47  
My name is andy lee

A multimedia application including video-conferencing

# Application Protocol

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## ■ URL

- Uniform resource locator
- <http://www.sharecourse.net/sharecourse/>

## ■ HTTP

- Hyper Text Transfer Protocol

## ■ TCP

- Transmission Control Protocol

## ■ 17 messages for one URL request

- 6 to find the IP (Internet Protocol) address
- 3 for connection establishment of TCP
- 4 for HTTP request and acknowledgement
  - ▶ Request: I got your request and I will send the data
  - ▶ Reply: Here is the data you requested; I got the data
- 4 messages for tearing down TCP connection

# Outline

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- Applications
- **Network Connectivity**
- Network Architecture
- Network Performance

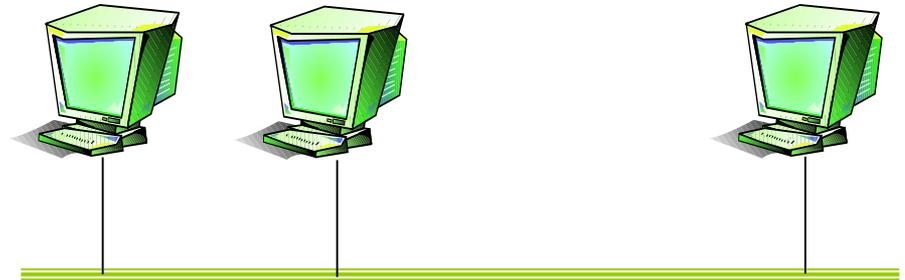
# Network Connectivity

## ■ Important terminologies

- Link
- Nodes
- Point-to-point
- Multiple access
- Switched Network
  - ▶ Circuit Switched
  - ▶ Packet Switched
- Packet, message
- Store-and-forward



Point-to-point

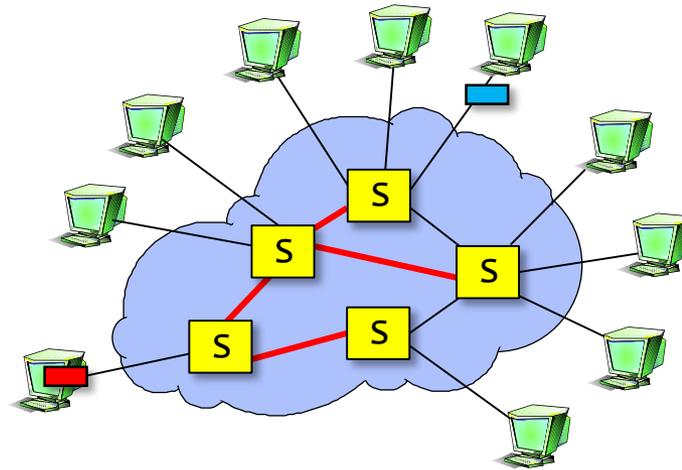


Multiple access

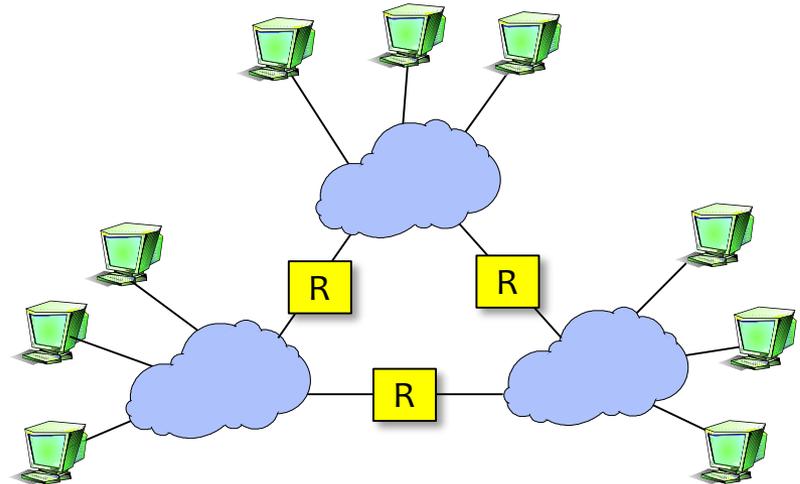
# Network Connectivity

## ■ Terminologies (contd.)

- Hosts
- Switches
- Spanning tree
- internetwork
- Router/gateway
- Host-to-host connectivity
- Address
- Routing
- Unicast/broadcast/multicast
- LAN (Local Area Networks)
- MAN (Metropolitan Area Networks)
- WAN (Wide Area Networks)

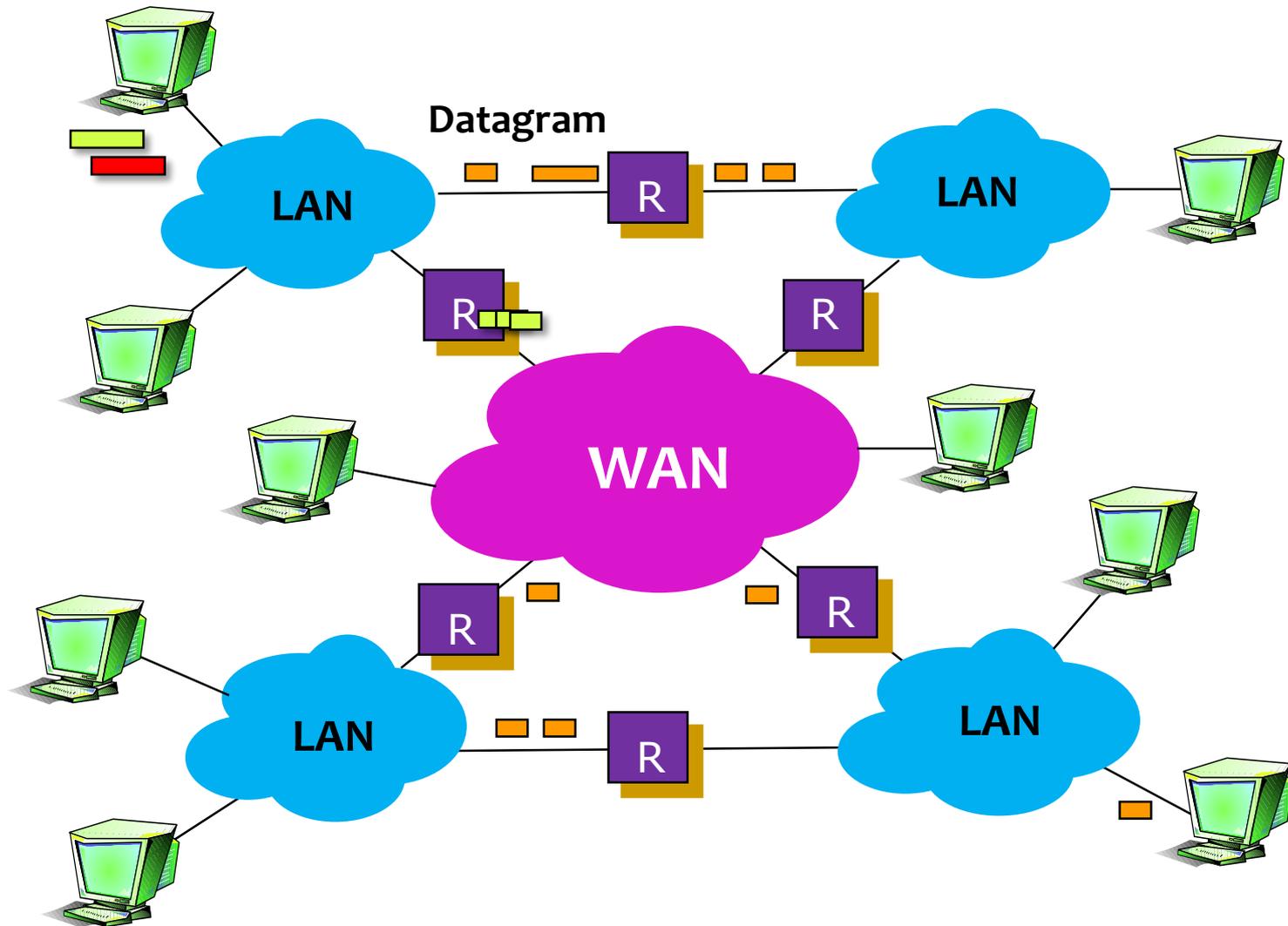


(a) A switched network



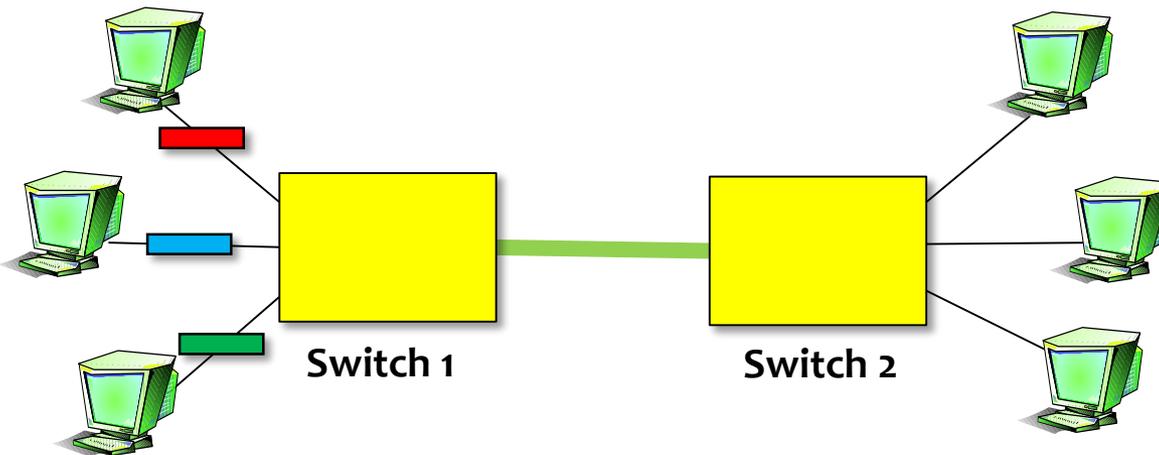
(b) Interconnection of networks

# How datagrams are delivered in an Internet ?



# Cost-Effective Resource Sharing

- Resource: links and nodes
- How to share a link ?
  - Multiplexing
  - De-multiplexing



Multiplexing multiple logical flows over a single physical link

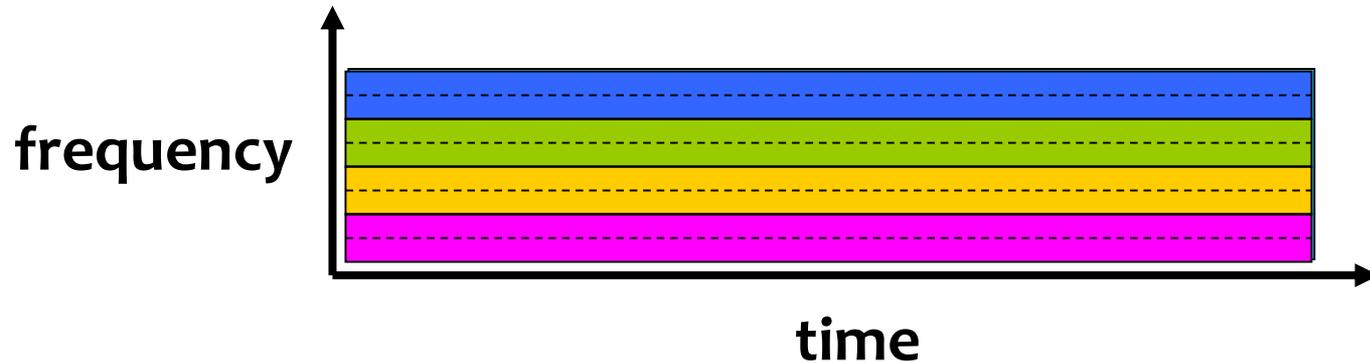
# Cost-Effective Resource Sharing

## ■ FDM: Frequency Division Multiplexing

FDM

Example:

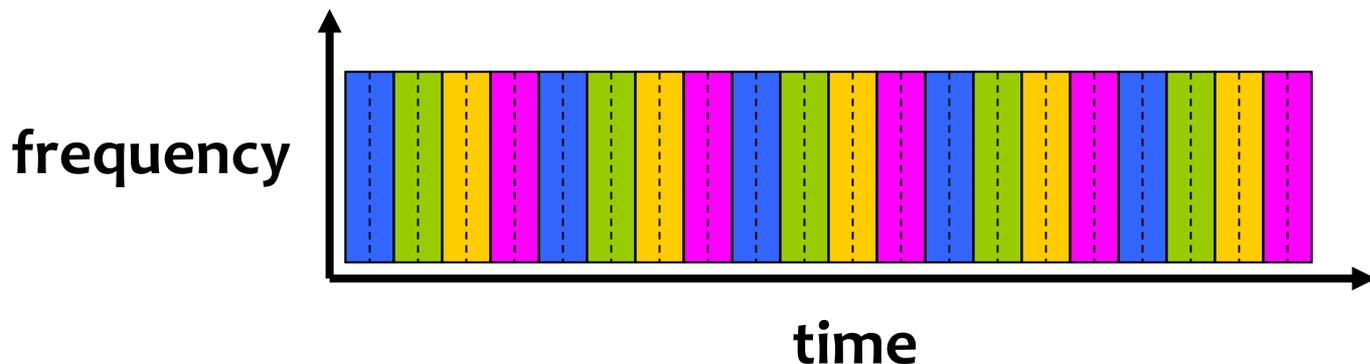
4 users



## ■ Synchronous Time-division Multiplexing (TDM)

- ▶ Time slots/data transmitted in predetermined slots

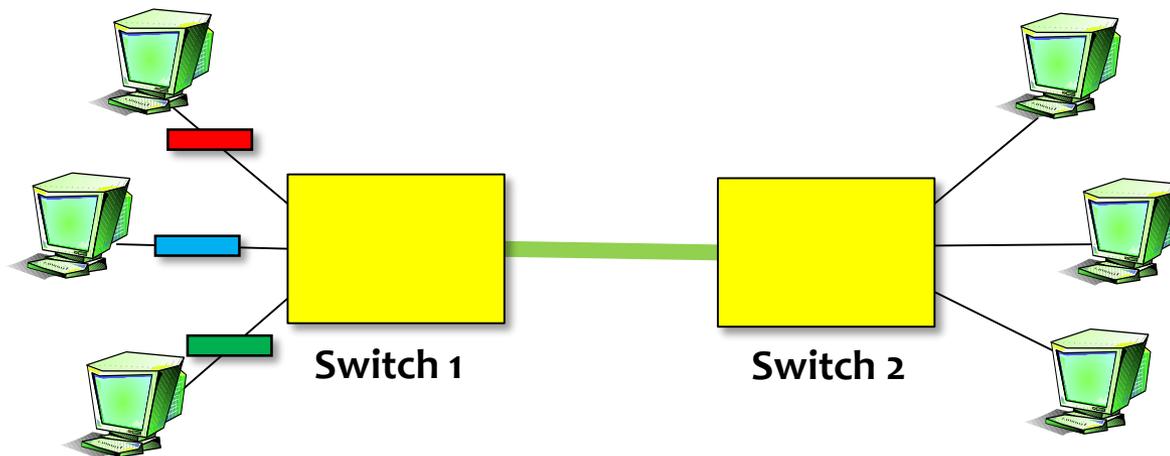
TDM



# Cost-Effective Resource Sharing

## ■ Statistical Multiplexing

- Data is transmitted based on demand of each flow.
- What is a **flow**?
- Packets vs. Messages
- FIFO, Round-Robin, Priorities (Quality-of-Service (QoS))
- Congested ?

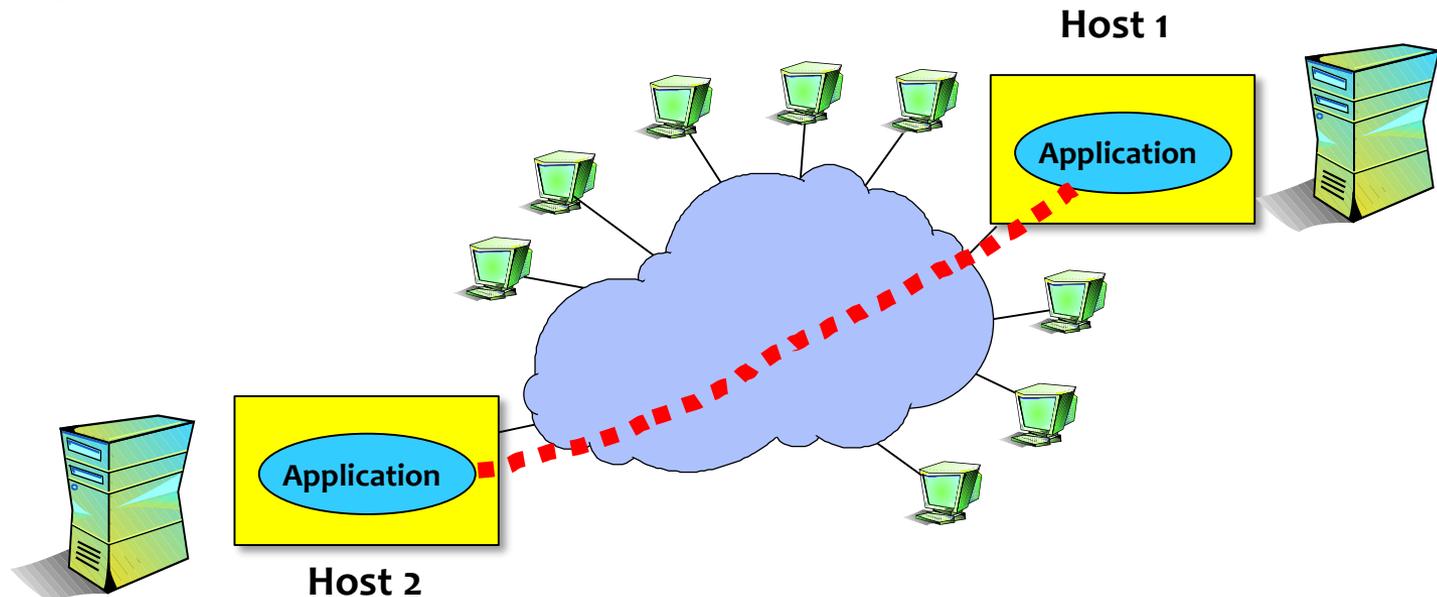


A switch multiplexing packets from multiple sources onto one shared link

# Logical Channels

## ■ Logical Channels

- Application-to-Application communication path or a pipe



Process communicating over an abstract channel

# Network Reliability

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- Network should **hide the errors**
- Bits are lost
  - Bit errors (1 to a 0, and vice versa)
  - Burst errors – several consecutive errors
- Packets are lost (Congestion)
- Links and Node failures
- Messages are delayed
- Messages are delivered out-of-order
- Third parties eavesdrop

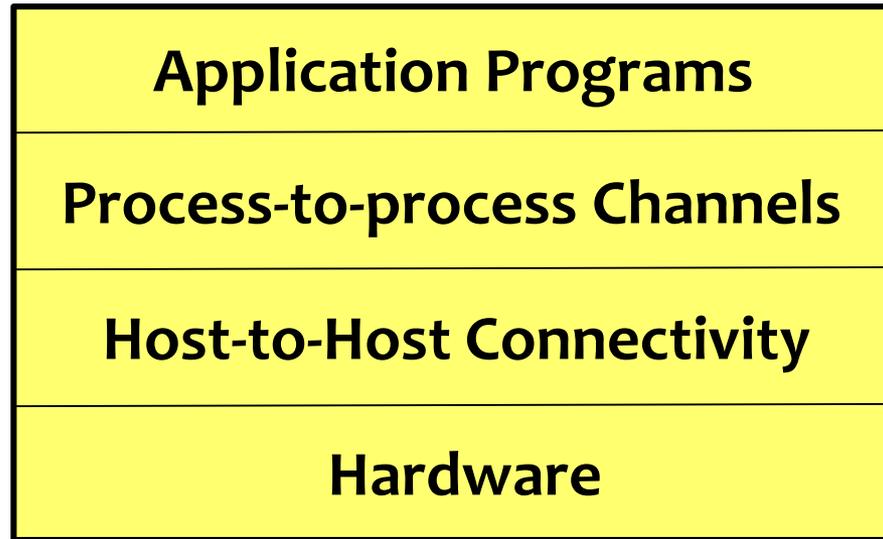
# Outline

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- Applications
- Network Connectivity
- **Network Architecture**
- Network Performance

# Network Architecture

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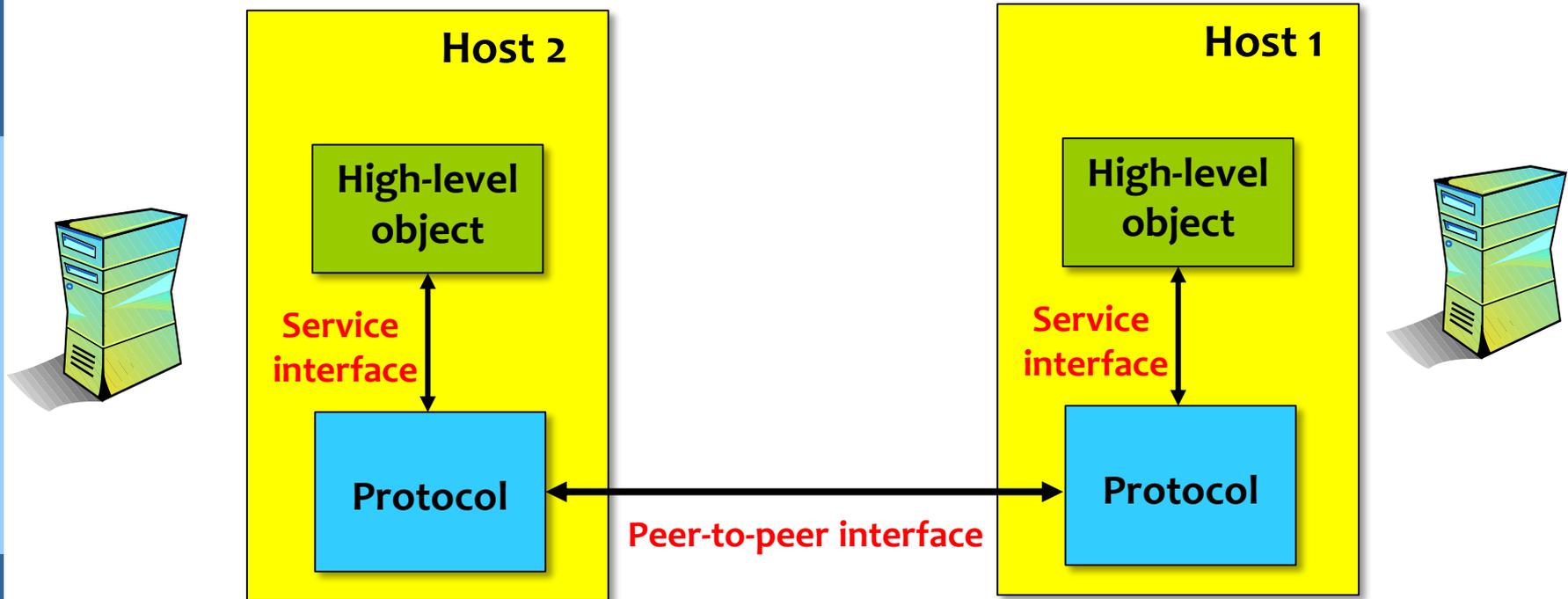
**Example of a layered network system**

# Protocols

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- **Protocol** defines the **interfaces** between
  - the layers in the same system and with
  - the layers of peer system
- Building blocks of a **network architecture**
- Each protocol object has two different interfaces
  - **Service interface**: operations on this protocol
  - **Peer-to-peer interface**: messages exchanged with peer

# Protocol Interfaces



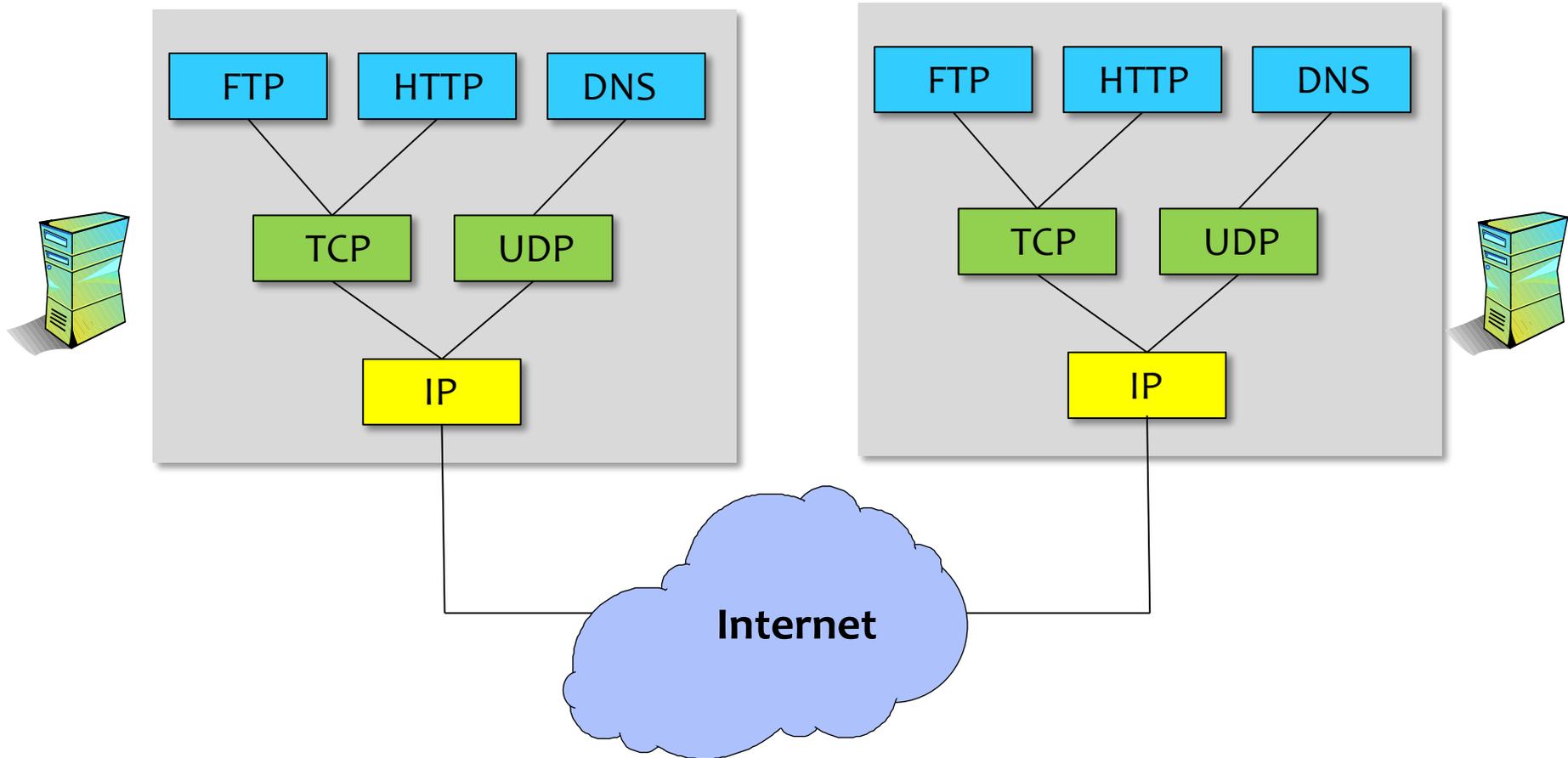
Service and Peer Interfaces for a protocol

# Protocols

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- **Protocol Specification:** pseudo-code, state transition diagram, message format
- **Interoperable:** when two or more protocols that implement the specification accurately
- **IETF: Internet Engineering Task Force**
  - Define Internet standard protocols

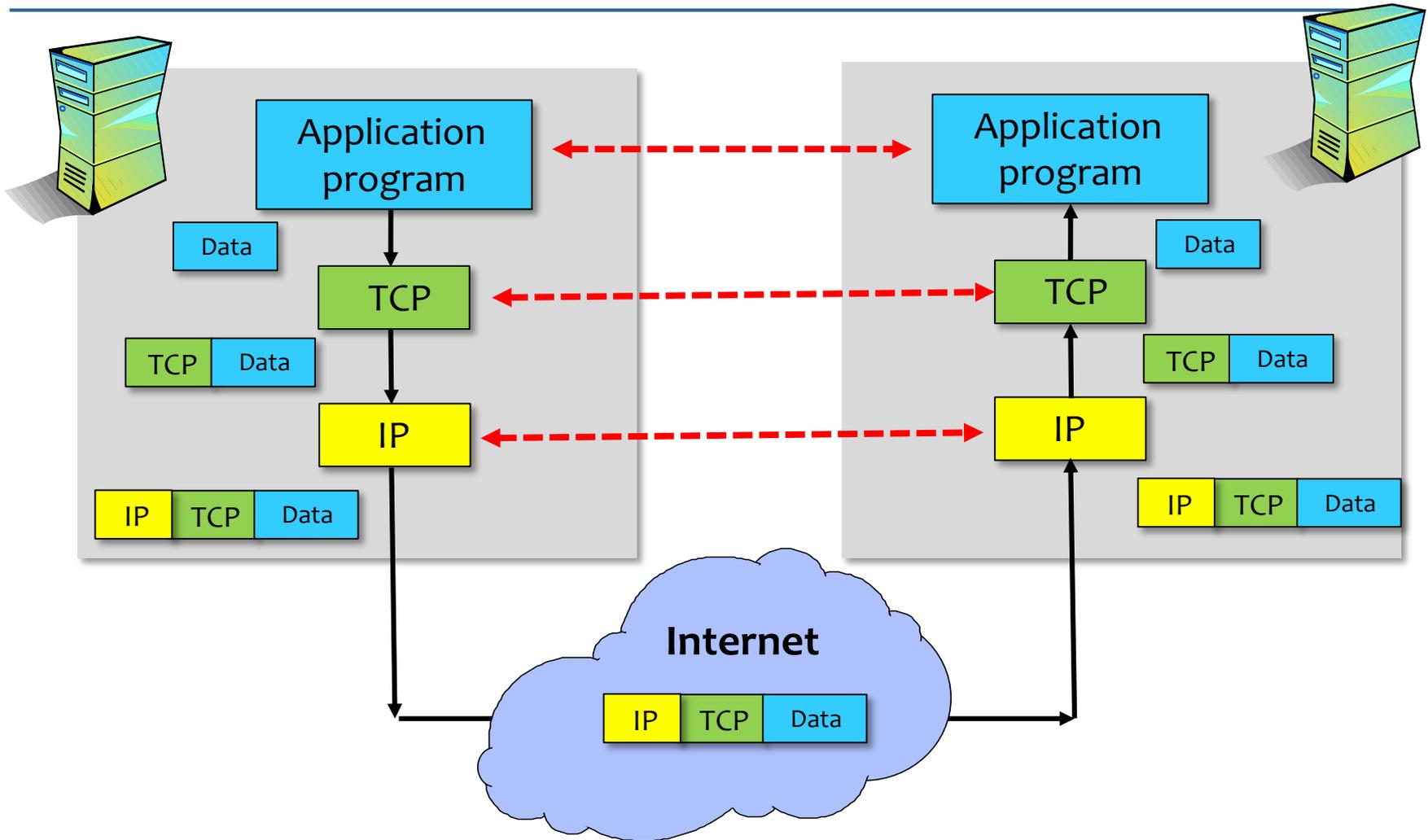
# Protocol Architecture



**Example of a protocol architecture**

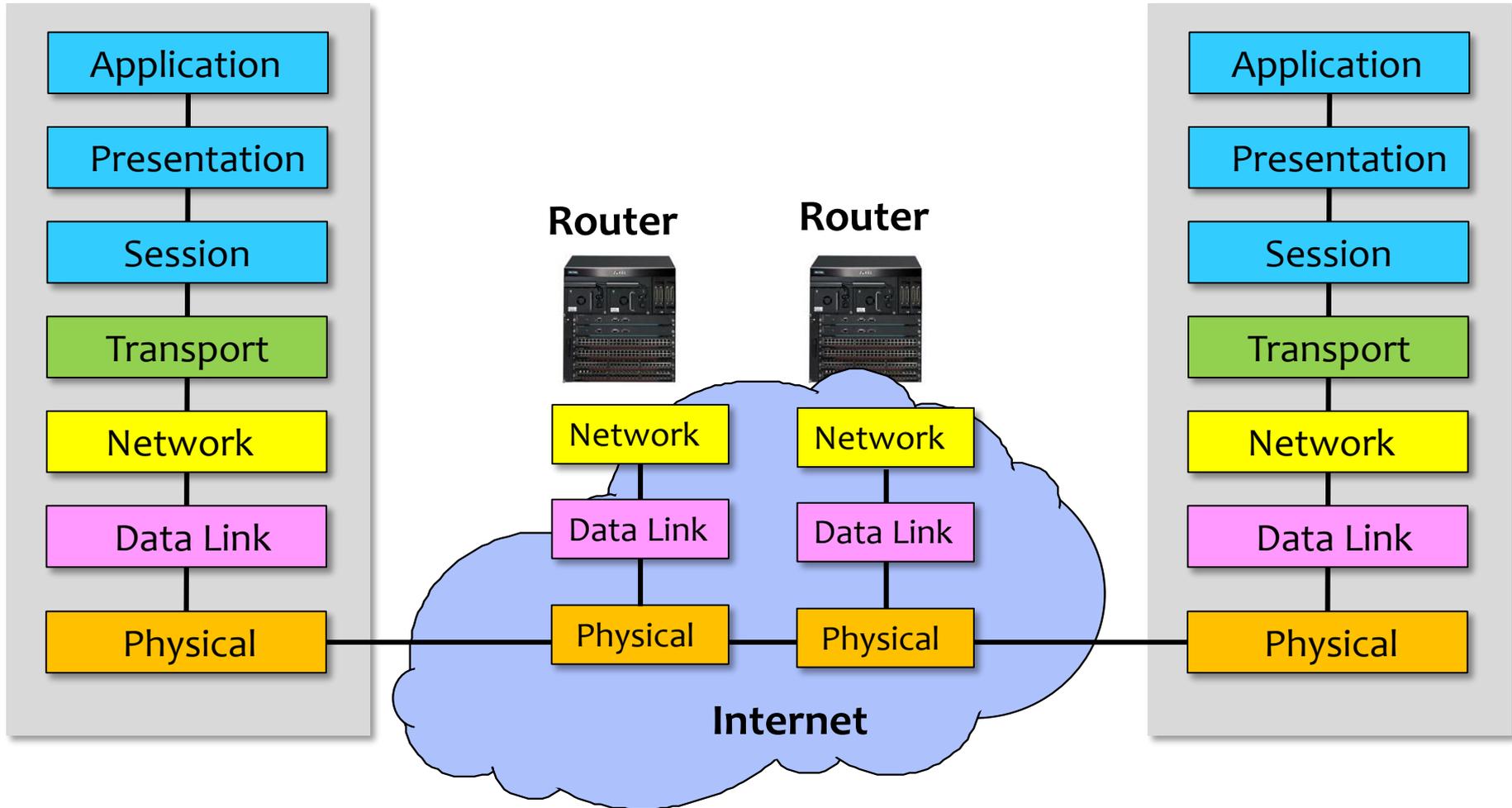
**nodes are the protocols and links the “depends-on” relation**

# Encapsulation



High-level messages are encapsulated inside of low-level messages

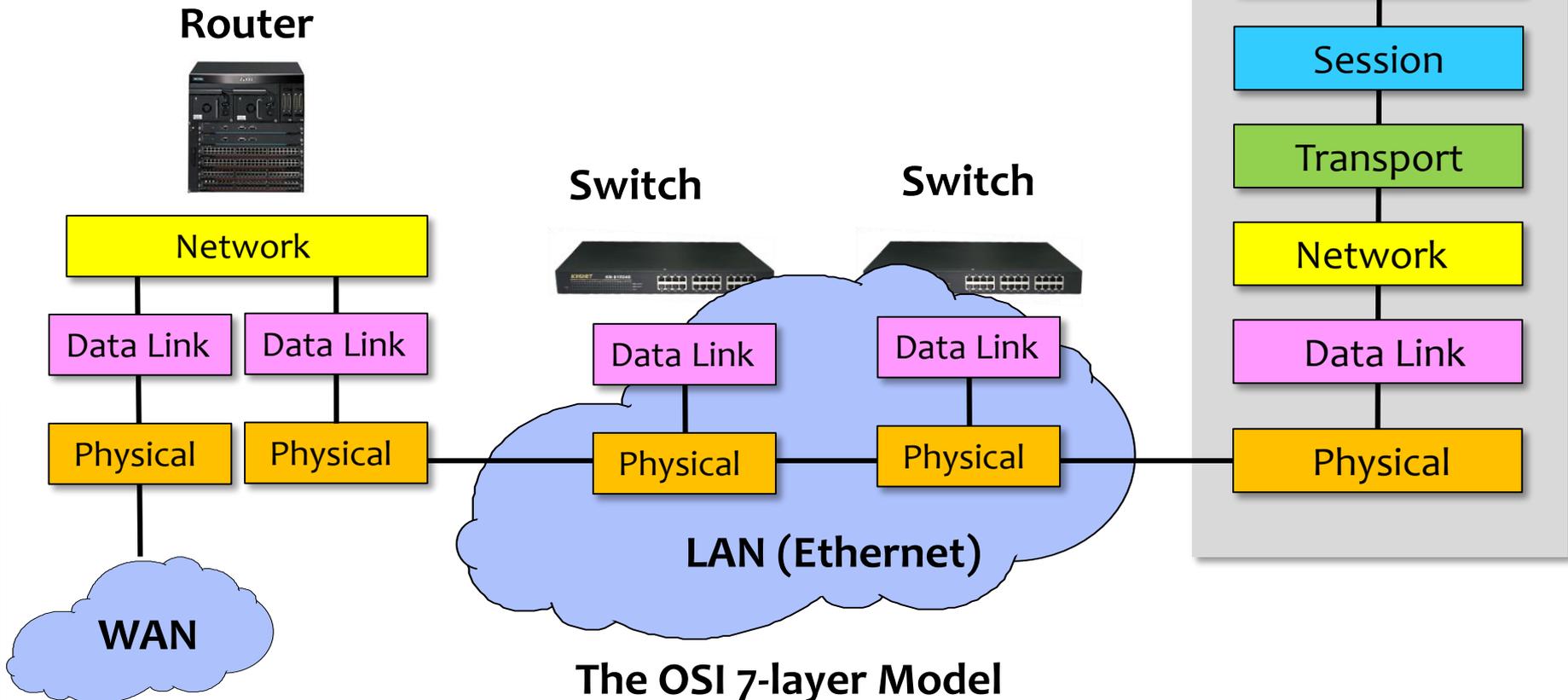
# OSI Architecture



The OSI 7-layer Model

OSI – Open Systems Interconnection

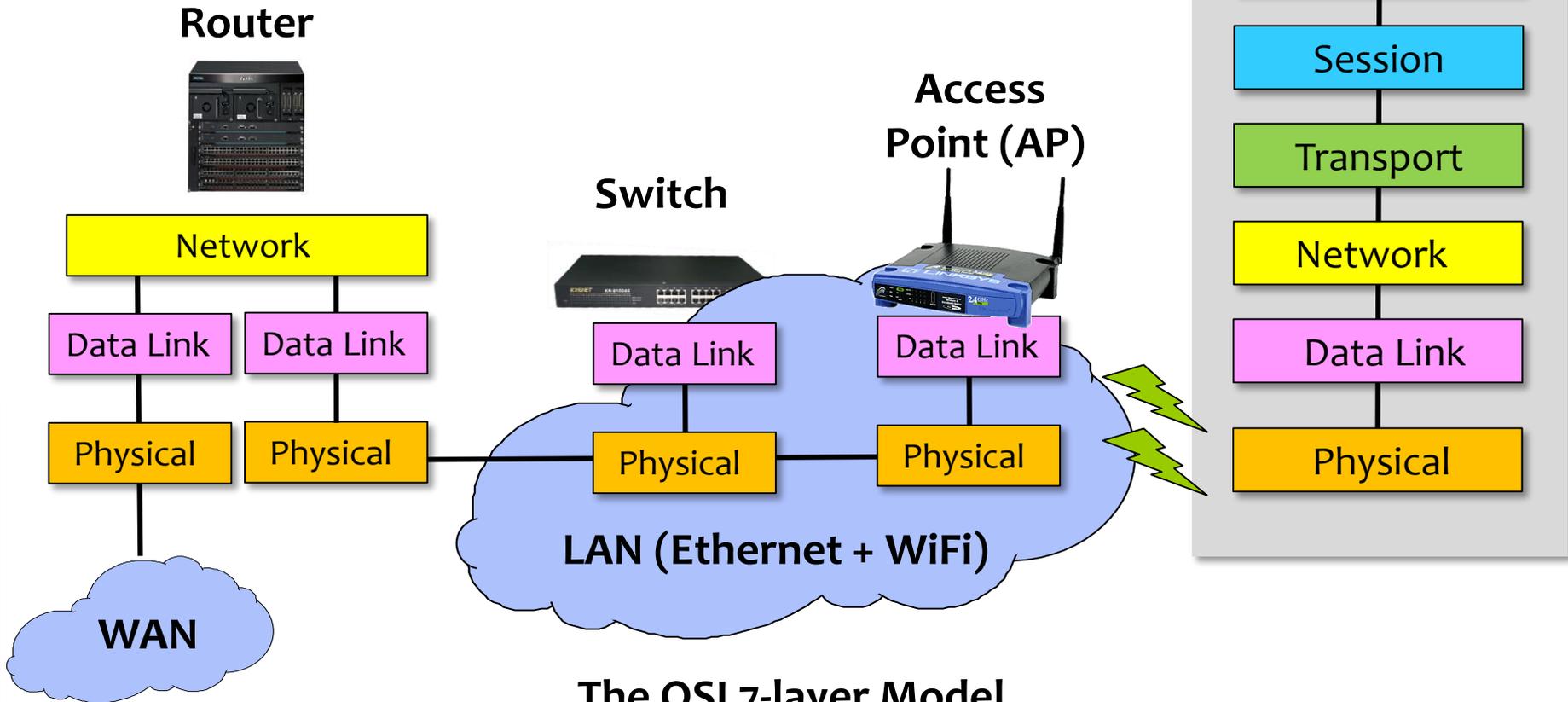
# OSI Architecture



The OSI 7-layer Model

OSI – Open Systems Interconnection

# OSI Architecture



The OSI 7-layer Model

OSI – Open Systems Interconnection

# Description of Layers

## ■ Physical Layer (如何將原始資料在 link 上傳輸)

- Handles the transmission of **raw bits over a communication link**

- ▶ **Coaxial cable**



- ▶ **Twisted pair**



- ▶ **Optical Fiber**

- ▶ **Air space (wireless radio channel)**



- Different Signal Coding schemes



# Description of Layers

## ■ Data Link Layer (如何將 frame 傳給直接相連的主機或設備)

- Collects a stream of bits into a *frame*
- How to transmit a frame to a directly connected host (destination) ?

- MAC (Media Access Control Protocol)

- ▶ CSMA/CD (IEEE 802.3 Ethernet)
- ▶ CSMA/CA (IEEE 802.11 Wireless LAN)

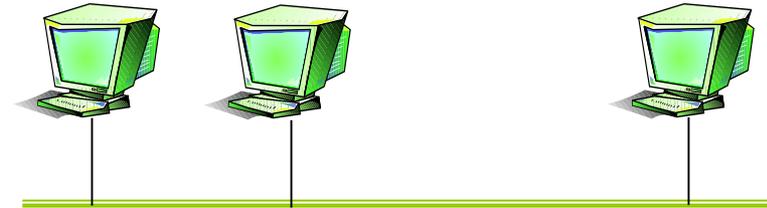
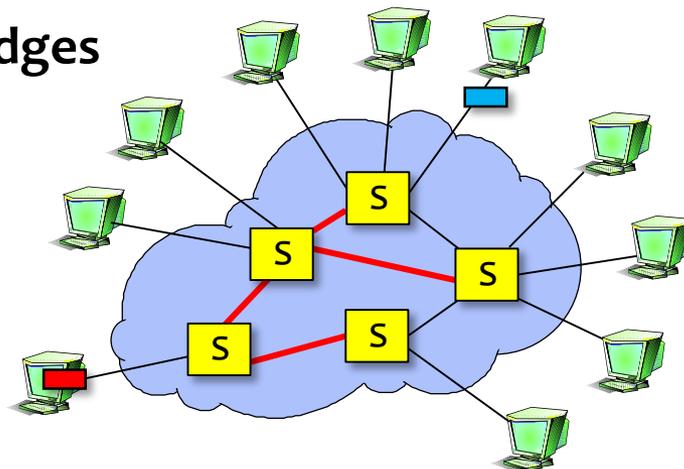


Point-to-point

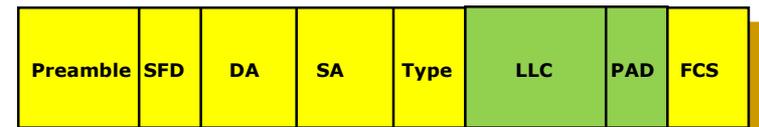
- Layer 2 devices

- Switches

- Bridges



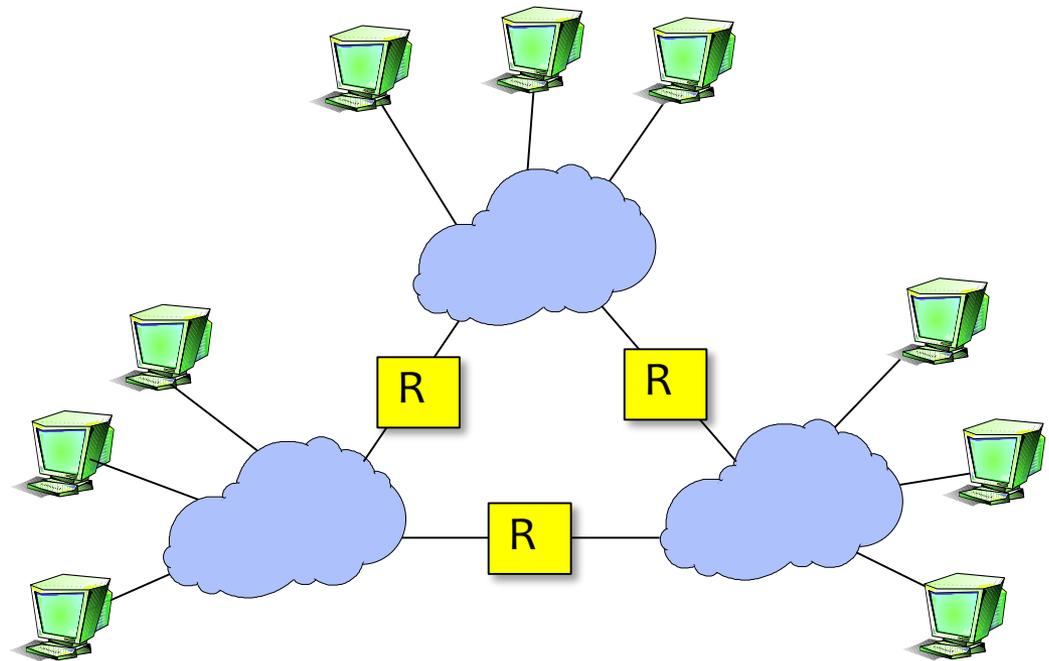
Multiple access



# Description of Layers

## ■ Network Layer (如何將封包透過 Internet 送給目的地主機)

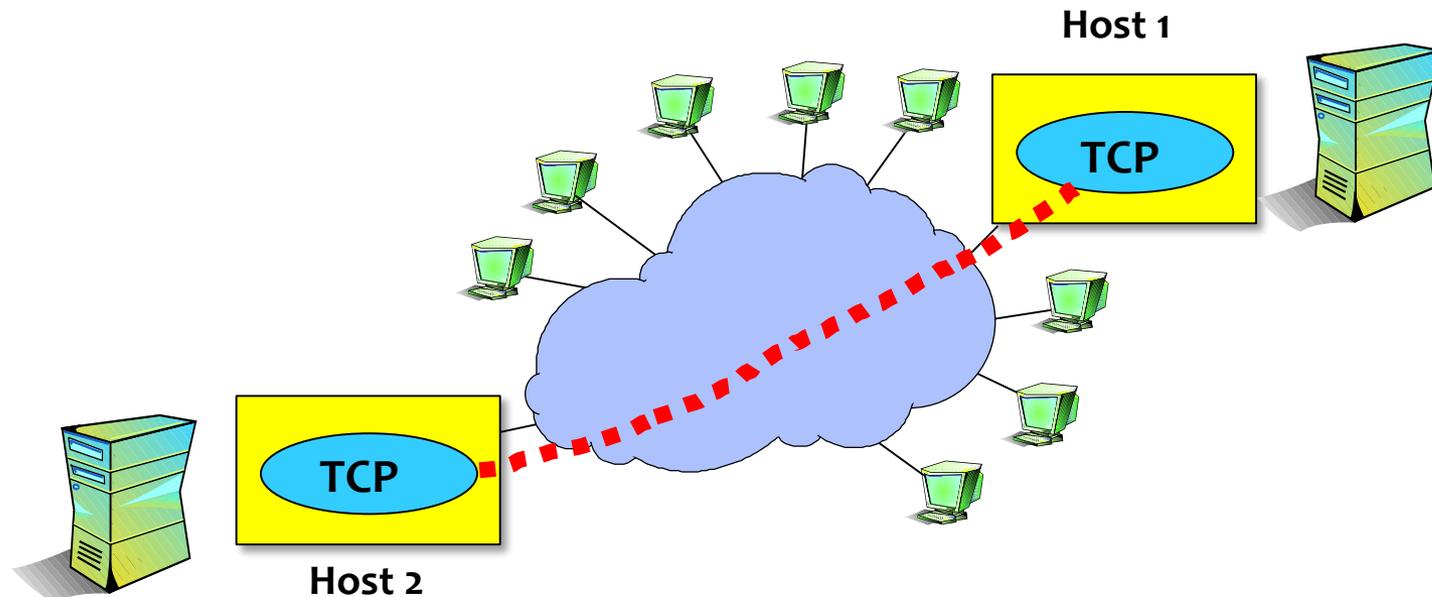
- How to transmit frames to a host via the Internet ?
- Handles **routing** among nodes within a packet-switched network
- Data exchanged between nodes in this layer is called a **packet**
- *IP protocol*
- *Routers*
- *Routing protocols*
  - ▶ *RIP*
  - ▶ *OSPF*
  - ▶ *BGP*
- *Routing Tables*



Interconnection of networks

# Description of Layers

- Transport Layer (提供不同主機 processes 之間的資料傳送)
  - Implements a **process-to-process** channel
  - Unit of data exchanges in this layer is called a **message**
  - **TCP (Transmission Control Protocol)** – *Reliable service*
  - **UDP (User Datagram Protocol)** – *Unreliable service*



# Description of Layers

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## ■ Session Layer

- Provides a name space that is used **to tie together the potentially different transport streams** that are part of a single application

## ■ Presentation Layer

- Concerned about the **format of data** exchanged between peers

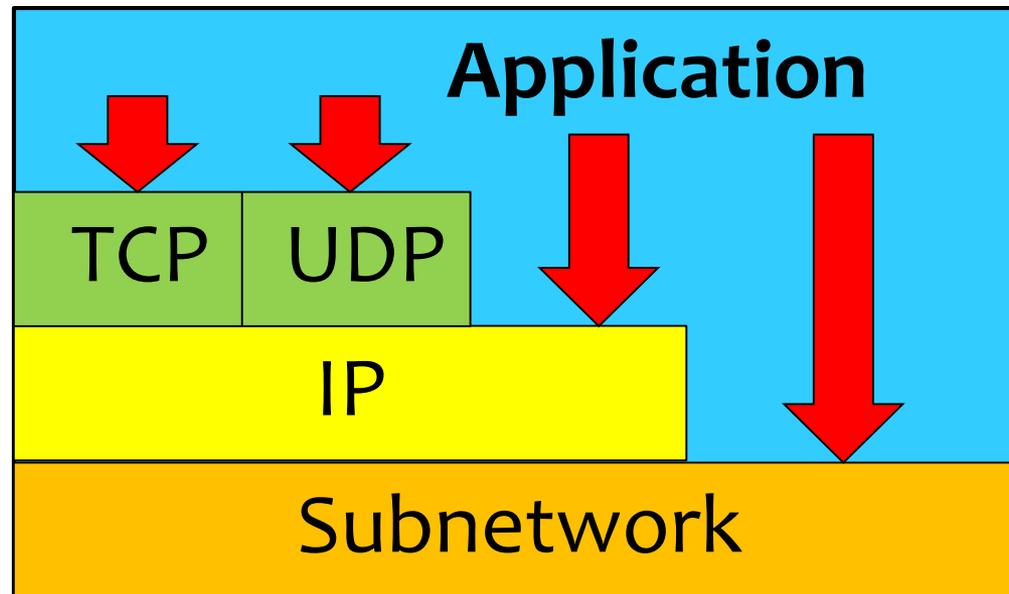
## ■ Application Layer

- Standardize common type of exchanges
- FTP/E-mail/DNS/HTTP/Browsers/FB, ....

**The transport layer and the higher layers typically run only on end-hosts and not on the intermediate switches and routers**

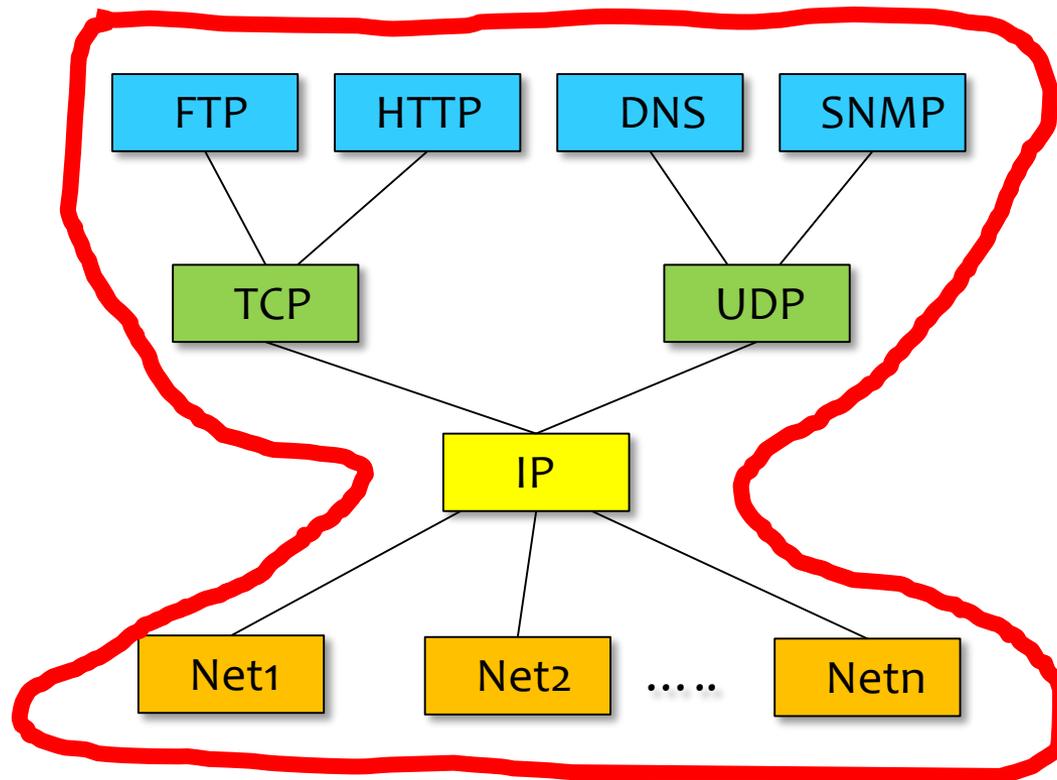
# Internet Architecture

- Defined by IETF
- Three main features
  - **Does not imply strict layering.** The application is **free to bypass** the defined transport layers and to directly use IP or other underlying networks



# Internet Architecture

- **An hour-glass shape** – wide at the top, narrow in the middle and wide at the bottom. IP serves as the focal point for the architecture



# Internet Architecture

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- In order for a new protocol to be officially included in the architecture, there needs to be both **a protocol specification** and **at least one** (and preferably two) **representative implementations** of the specification

# Outline

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- **Network Performance**

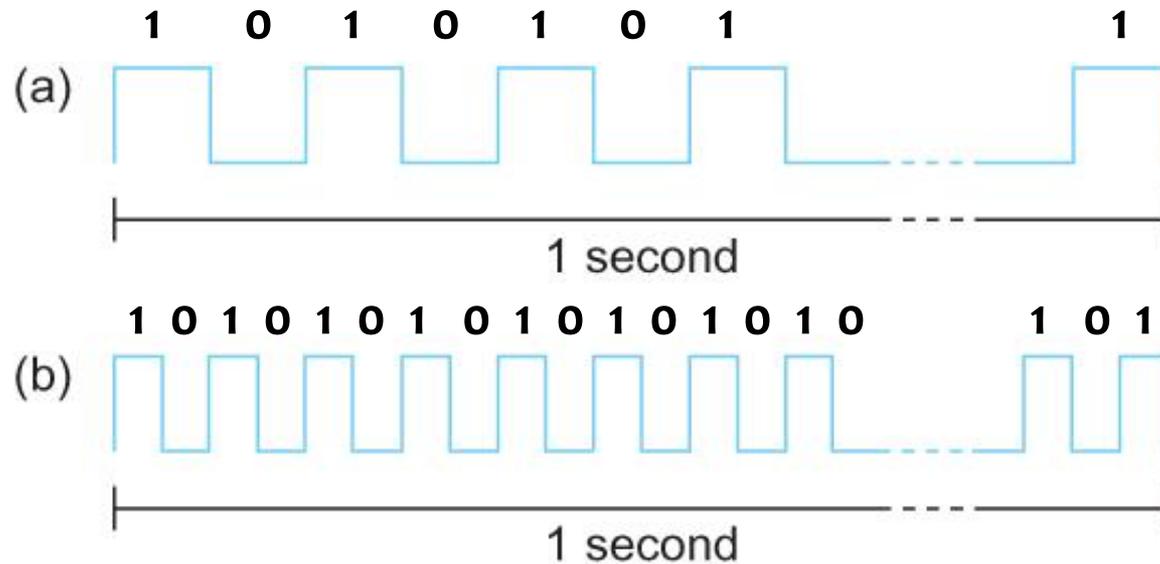
# Network Performance

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## ■ Bandwidth

- Width of the **frequency band**
  - **Number of bits per second** that can be transmitted over a communication link
- 
- 1 Mbps:  $1 \times 10^6$  bits/second
  - $1 \times 10^{-6}$  seconds to transmit each bit or imagine that a timeline, now each bit occupies 1 micro second space.
  - On a 2 Mbps link the width is 0.5 micro second.
  - Smaller the width more will be transmission per unit time.

# Bandwidth



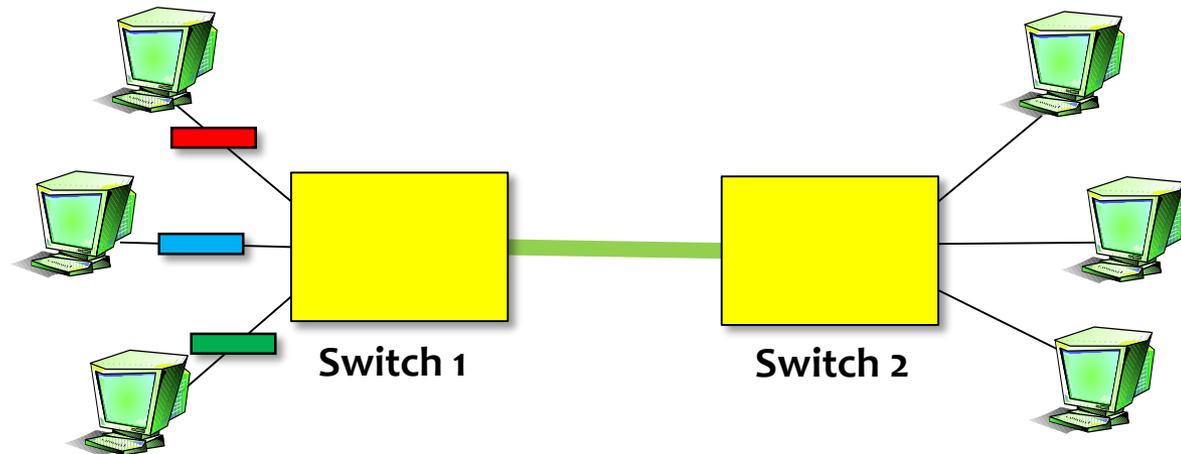
**Bits transmitted at a particular bandwidth can be regarded as having some width:**

**(a) bits transmitted at 1Mbps (each bit 1  $\mu$ s wide);**

**(b) bits transmitted at 2Mbps (each bit 0.5  $\mu$ s wide).**

# Network Performance

- **Latency** = Propagation time + transmission time + queuing time
- **Propagation time** = distance/speed of light
- **Transmission time** = size/bandwidth



- One bit transmission => propagation is important (短資料很快就送完, 但需要長時間才能傳到對方, 資料已送完, 但前導資料還未到達對方)
  - Propagation time >> transmission time
- Large bytes transmission => bandwidth is important (長資料很慢才能送完, 未送完前, 前導資料已到對方)
  - Transmission time >> propagation time

# Delay X Bandwidth

- The channel between a pair of processes can be viewed as a pipe
- **Latency (delay): length** of the pipe
- **Bandwidth: width** of the pipe
- Delay x Bandwidth means how many data can be stored in the pipe
- For example, delay of 80 ms and bandwidth of 100 Mbps
  - ⇒  $80 \times 10^{-3}$  seconds x  $100 \times 10^6$  bits/second
  - ⇒  $8 \times 10^6$  bits = 8 M bits = 1 MB data.



Network as a pipe

# Delay X Bandwidth

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- **Relative importance of bandwidth and latency depends on application**
  - **For large file transfer, bandwidth is critical**
  - **For small messages (HTTP, NFS, etc.), latency is critical**
  - **Variance in latency (**jitter**) can also affect some applications (e.g., audio/video conferencing)**

# Delay X Bandwidth

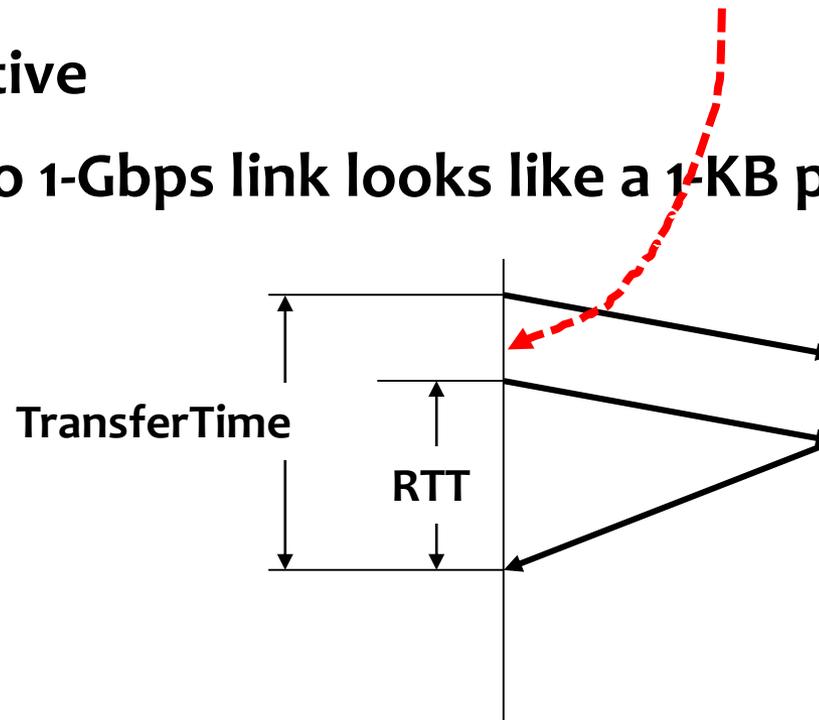
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- if the sender keeps the pipe full, **delay x bandwidth** is the number of bits the sender must transmit before the first bit arrives at the receiver
- Takes another one-way latency to receive a response from the receiver
- The sender will not fully utilize the network if the sender does not fill the pipe
  - ▶ send a whole delay x bandwidth product's worth of data before it stops to wait for a signal
  - ▶ 在停下來等對方回應之前應該要傳送 delay x bandwidth 的資料量

# Throughput

## ■ Infinite bandwidth

- RTT (Round Trip Time) dominates
- **Throughput** =  $\text{TransferSize} / \text{TransferTime}$
- $\text{TransferTime} = \text{RTT} + \text{TransferSize}/\text{Bandwidth}$
- Its all relative
- 1-MB file to 1-Gbps link looks like a 1-KB packet to 1-Mbps link



# Summary

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- A **layered architecture** for computer network
  - Physical Layer
  - Data Link Layer
  - Network Layer
  - Transport Layer
  - Session layer / Presentation Layer /Application layer
- **Two performance metrics** used to analyze the performance of computer networks
  - Bandwidth
  - Delay